iTAP: In-network Traffic Analysis Prevention using Software-Defined Networks



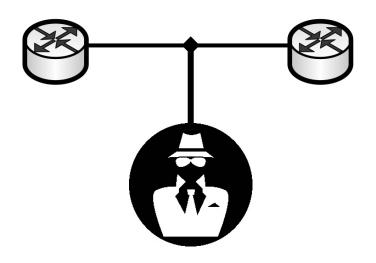
Roland Meier, David Gugelmann, Laurent Vanbever

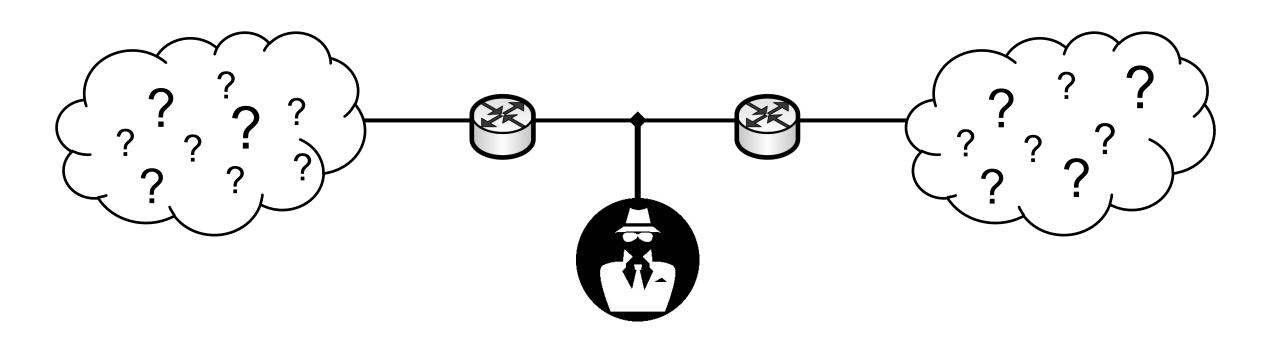
https://itap.ethz.ch

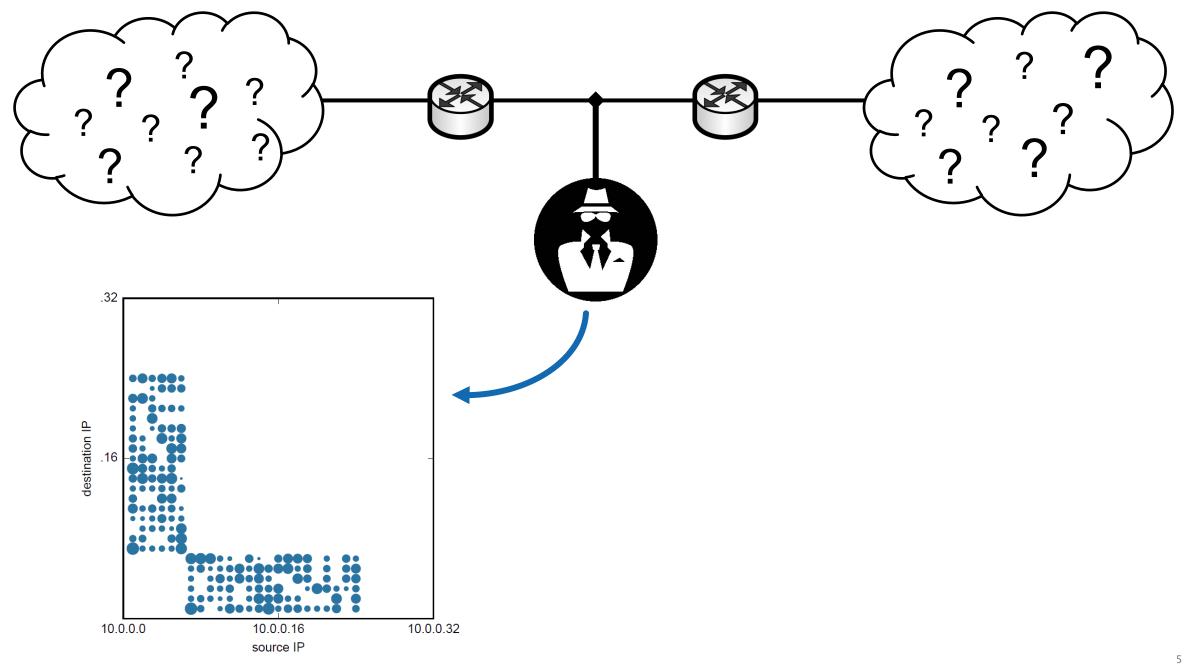
SDN Switzerland, 8th SDN Workshop. Zürich, CH (June 2017).

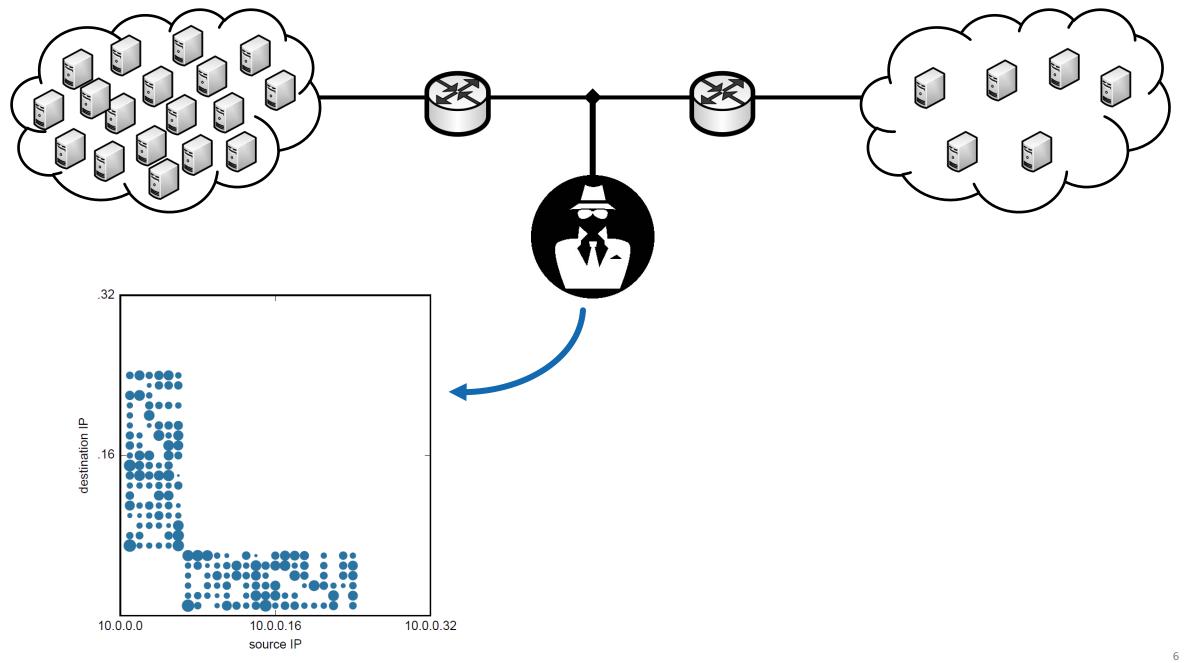


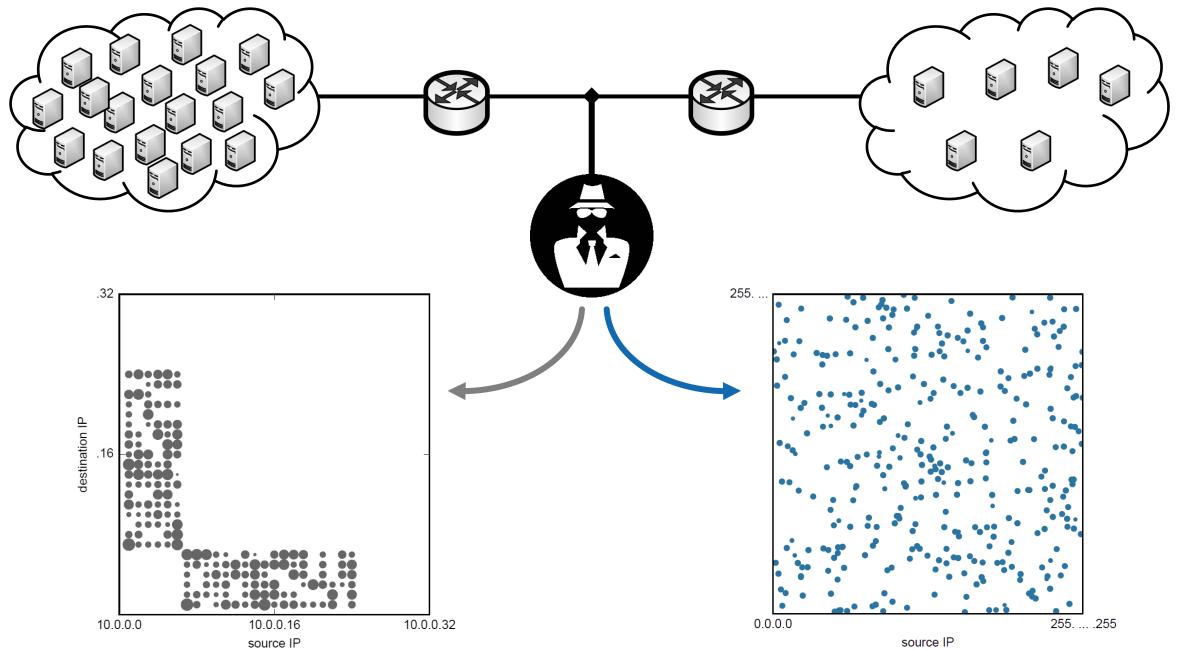


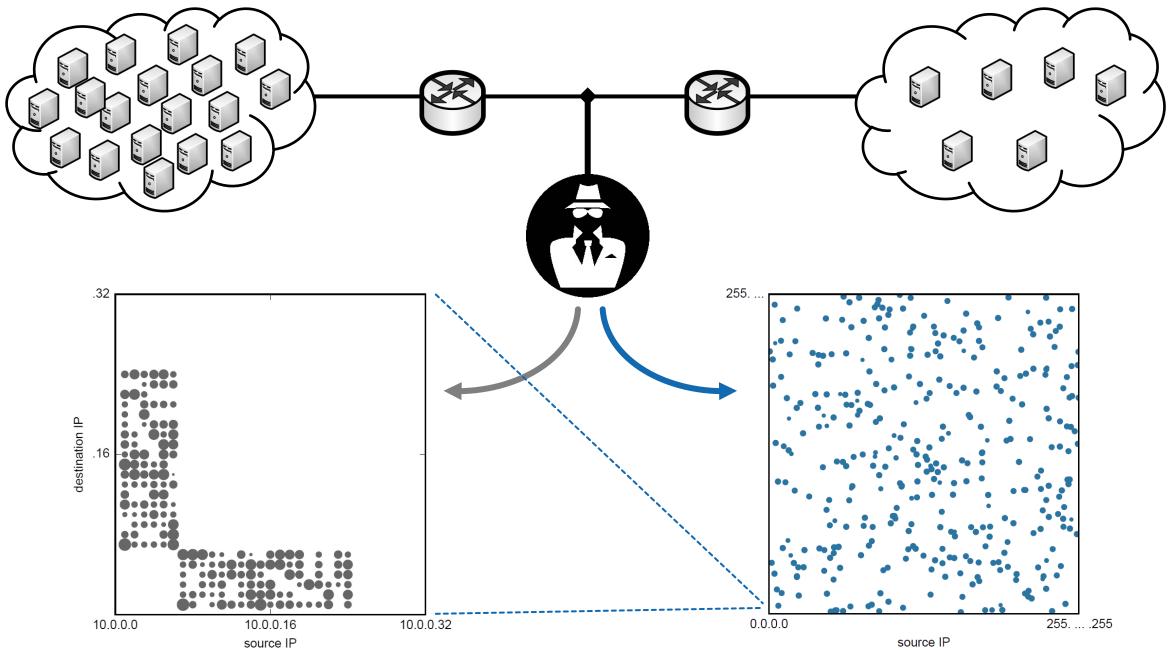


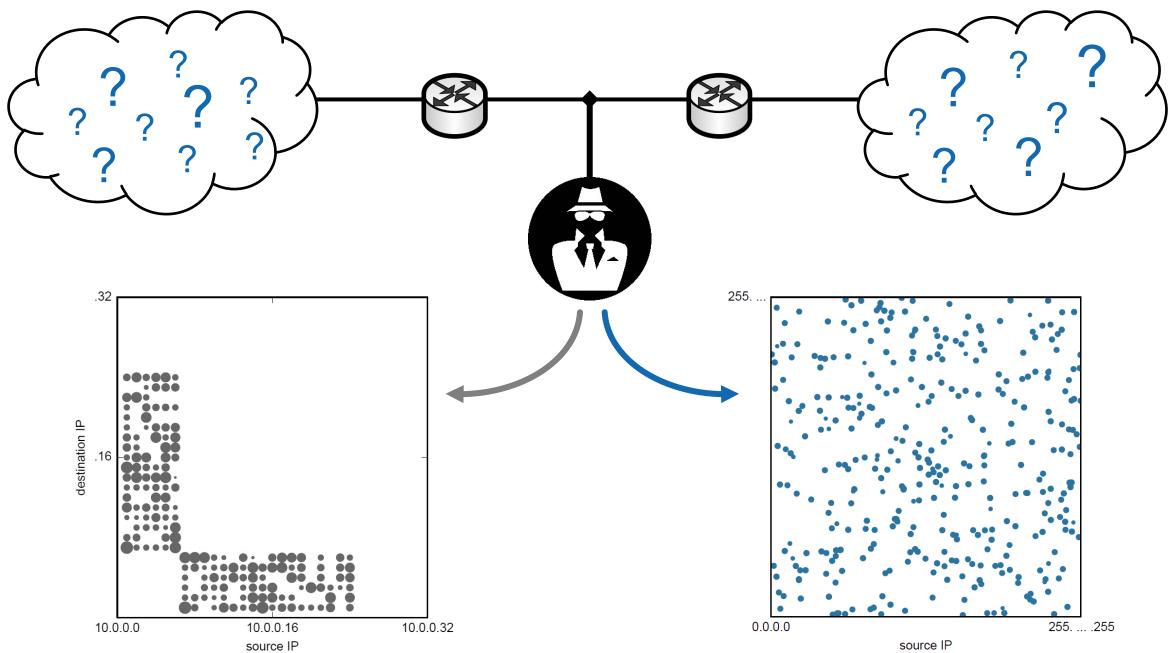


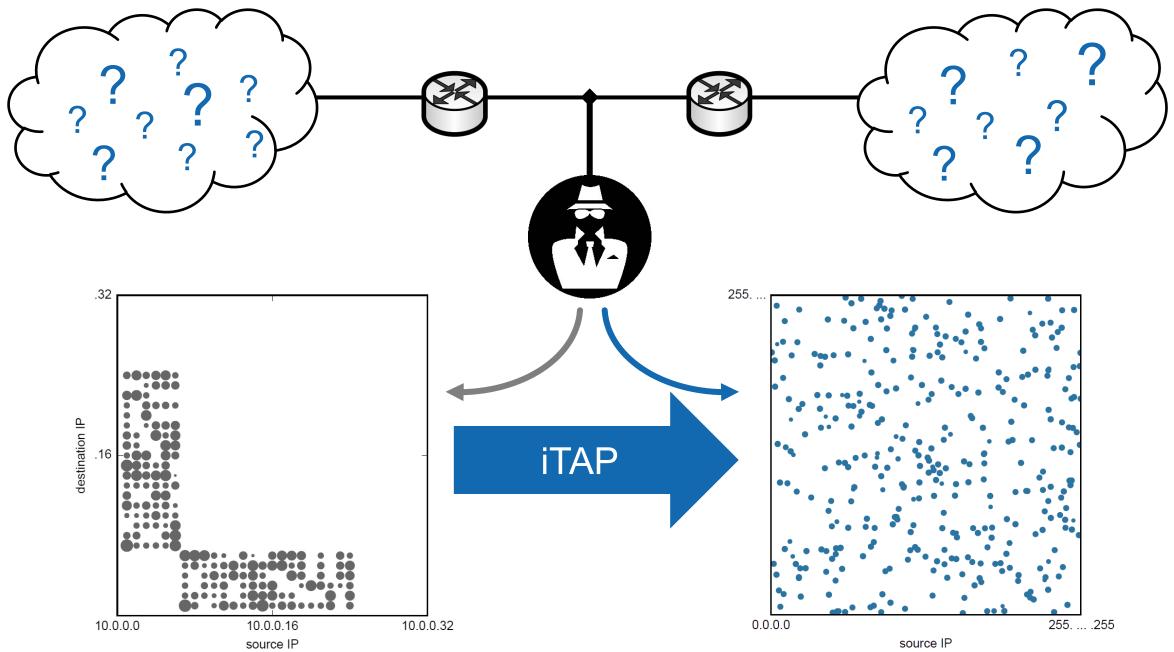












N.S.A. May Have Hit Internet Companies at a Weak Spot

The Internet companies' data centers are locked down with full-time security [...]. But between the data centers [...] information was unencrypted and an easier target for government intercept

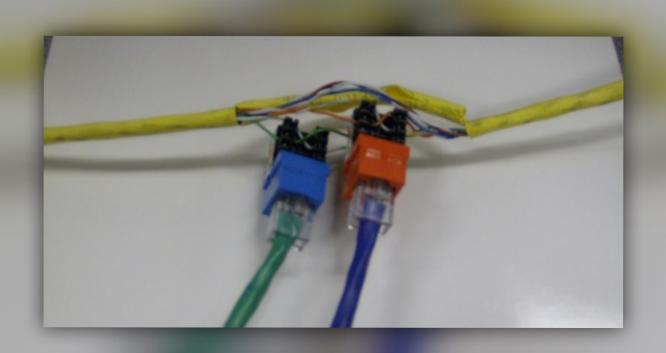
efforts, according to three people with knowledge of Google's and Yahoo's systems who spoke on the condition of

- The New York Times, Nov. 25, 2013 anonymity. MAtlantic SUBSCRIBE SEARCH MENU≡



Tapping

GLOBAL





Existing solutions

Do not protect communicating parties

[SSL/TLS, IPsec Transport, MACsec]

Require modifications at end-hosts or additional middleboxes [APOD, CONTRA]

Do not support partial deployment or have scalability problems [MACsec, PHEAR]

More references provided in the paper

In-network Traffic Analysis Prevention using Software-Defined Networks

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In-network Traffic Analysis Prevention using Software-Defined Networks

In-network Traffic Analysis Prevention using Software-Defined Networks

Communication anonymity who is communicating with whom?

In-network Traffic Analysis Prevention using Software-Defined Networks

- Communication anonymity who is communicating with whom?
- Volume anonymity how much traffic flows between X and Y?

In-network Traffic Analysis Prevention using Software-Defined Networks

- Communication anonymity who is communicating with whom?
- Volume anonymity how much traffic flows between X and Y?
- Topology anonymity how many hosts are in the network?

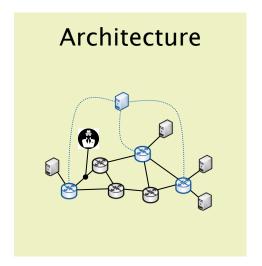
In-network Traffic Analysis Prevention using Software-Defined Networks

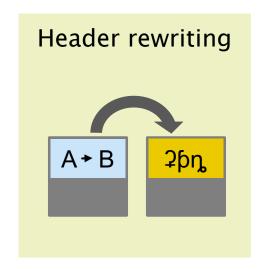
No modifications at end-hosts

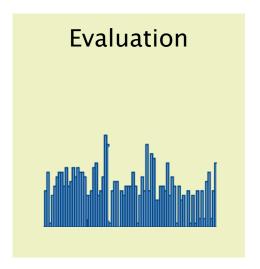
In-network Traffic Analysis Prevention using Software-Defined Networks

- Central controller
- Rewriting capabilities of switches

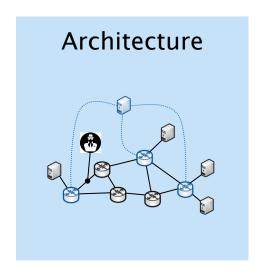


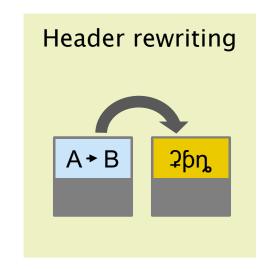


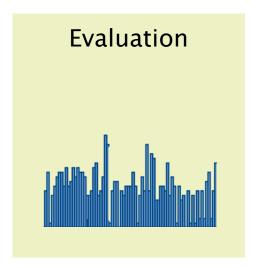


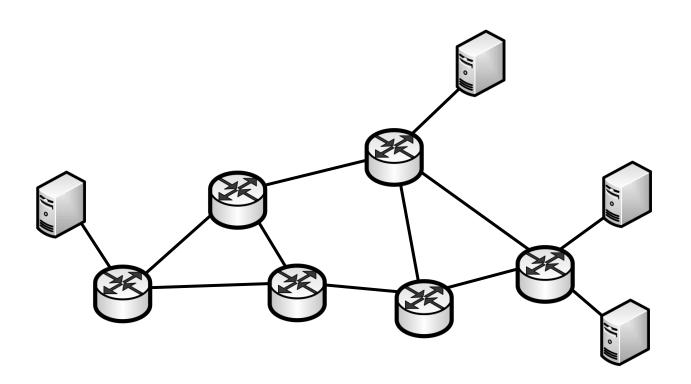


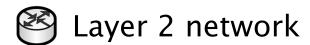


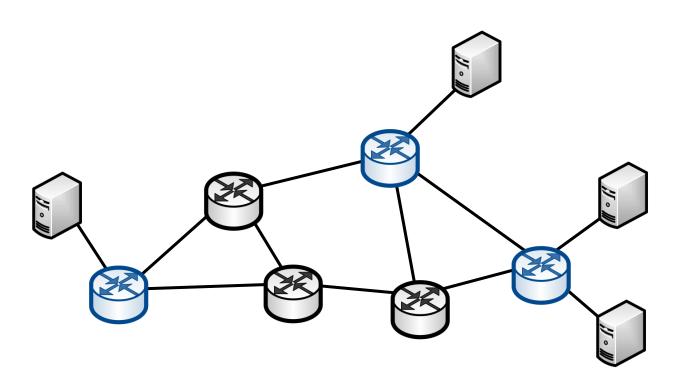




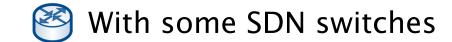


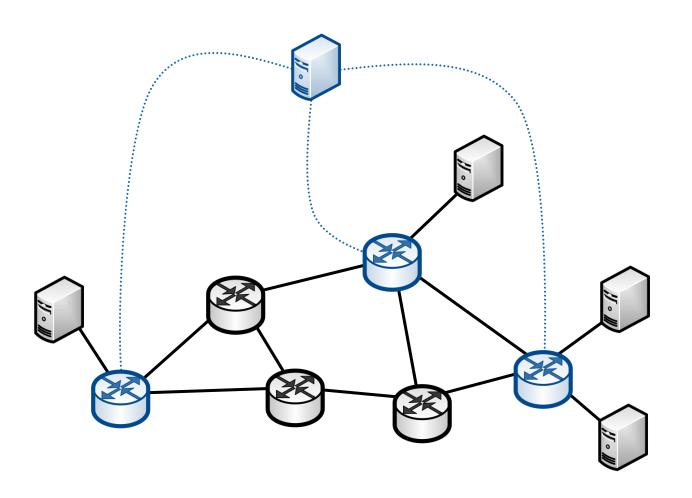








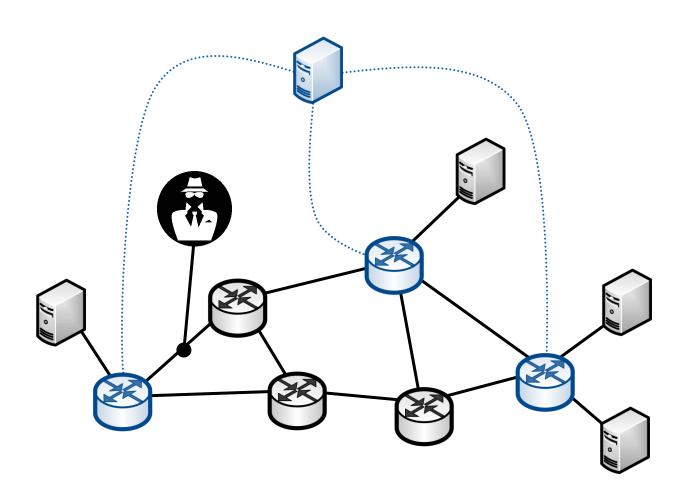








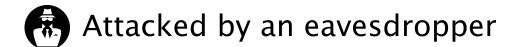


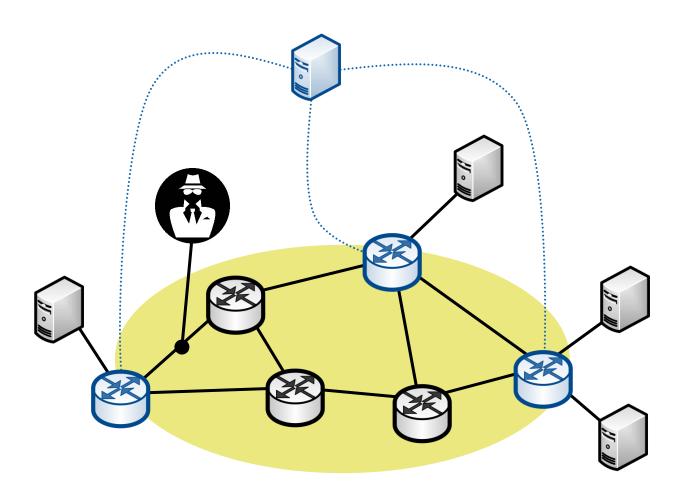














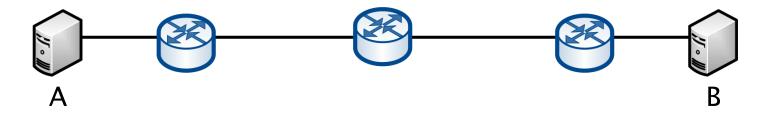




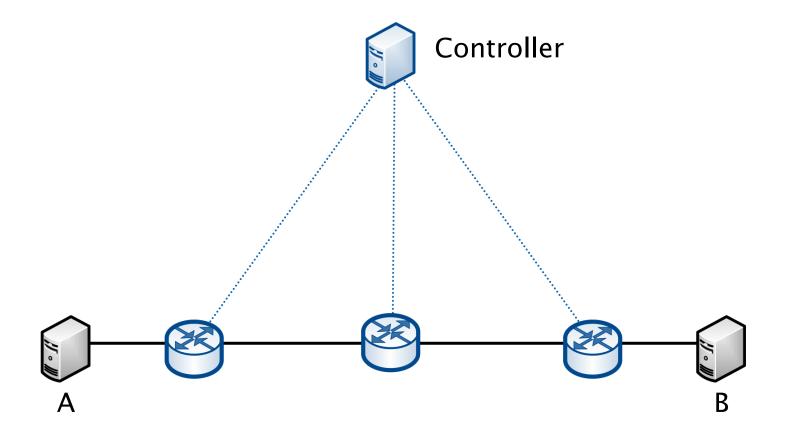
Attacked by an eavesdropper

Protected by iTAP

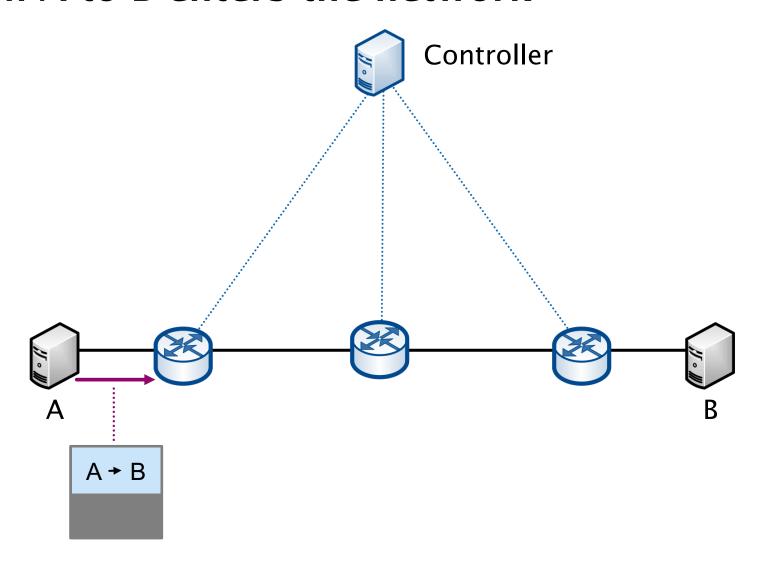
Example



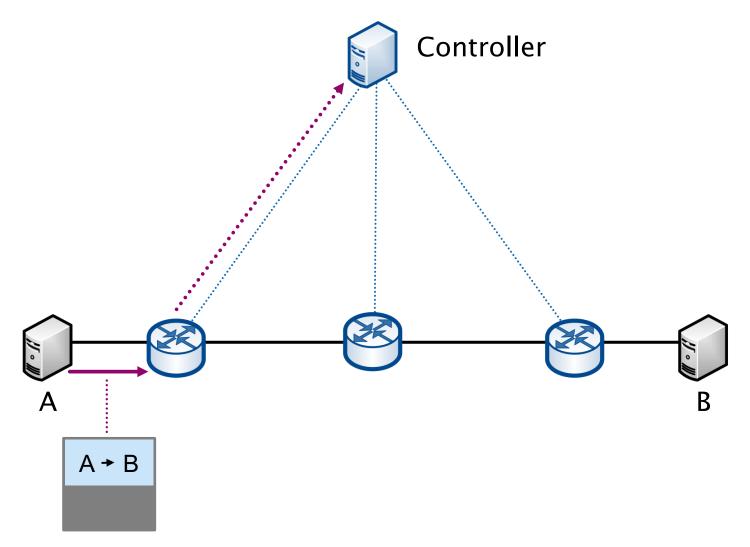
Example



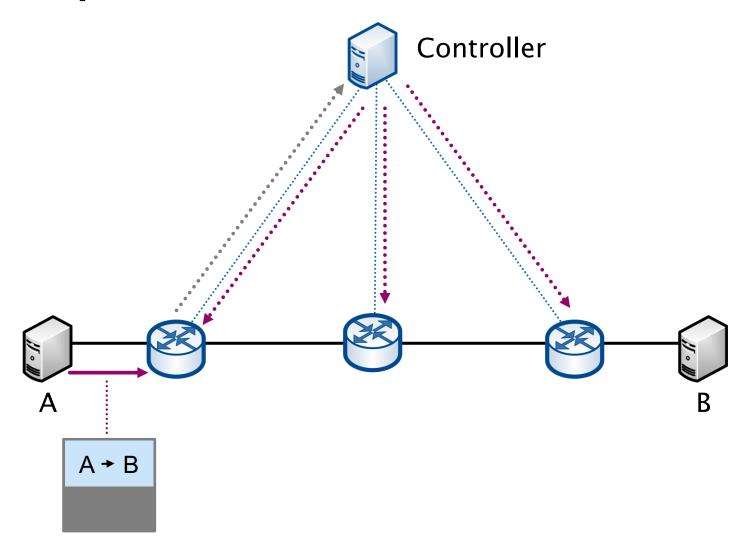
Packet from A to B enters the network



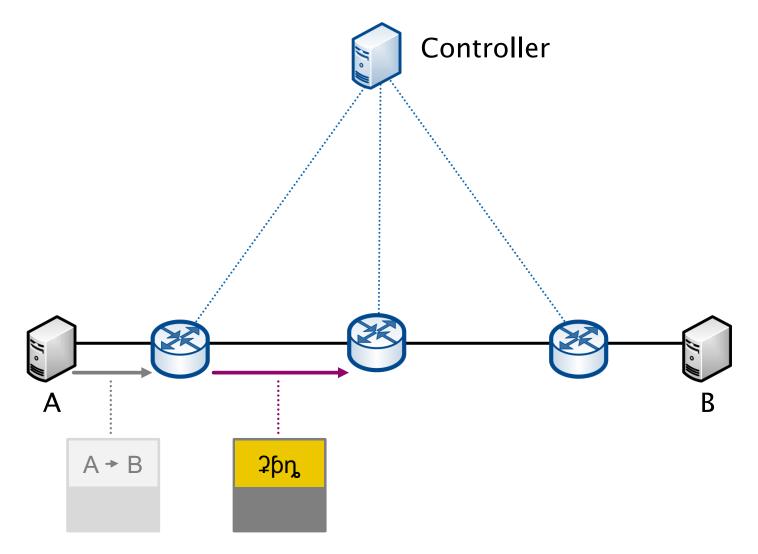
Ingress switch notifies controller



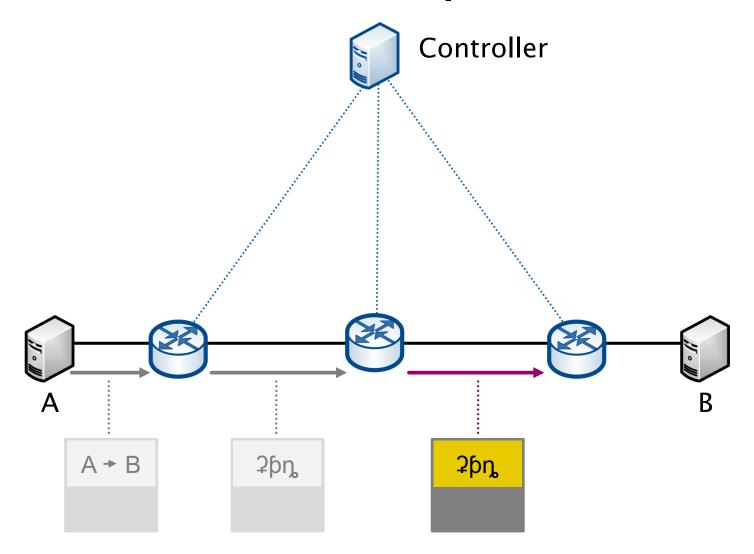
Controller computes & installs flow rules



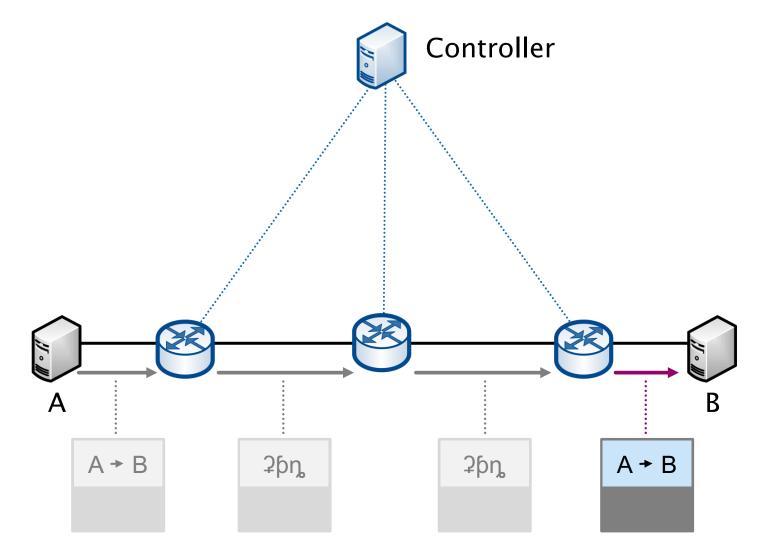
Ingress switch obfuscates source and destination



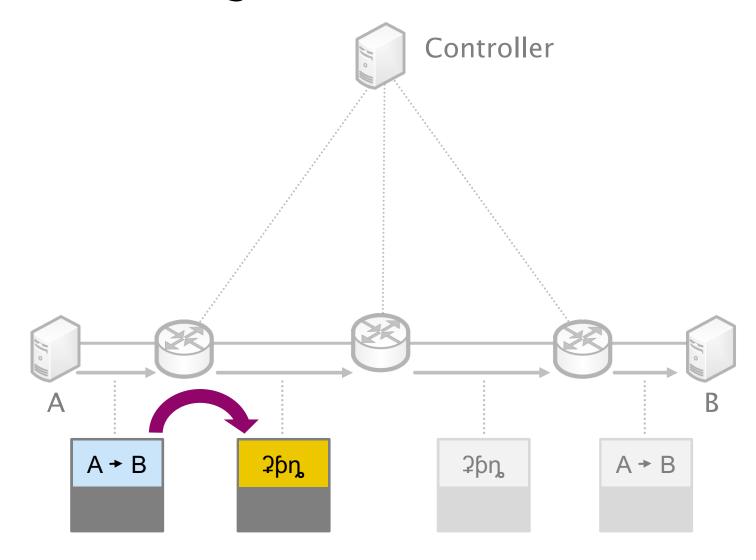
Core switch forwards obfuscated packet



Egress switch de-obfuscates source and destination

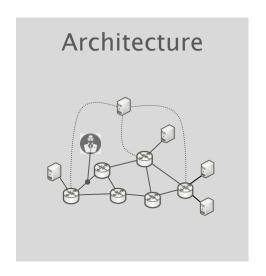


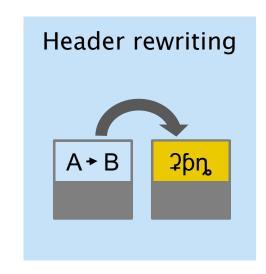
How does the rewriting work?

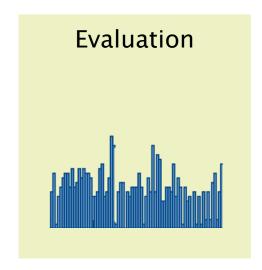


iTAP









Rewriting packet headers



Trade-off between anonymity and scalability



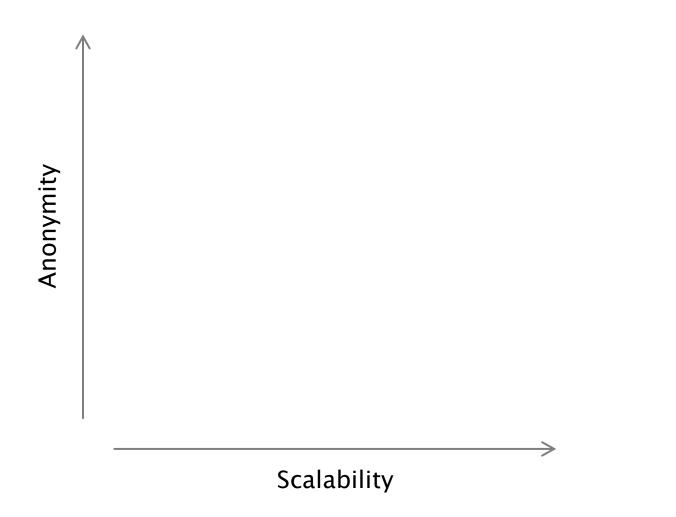
iTAP approach: Mixing per-host IDs and random bits

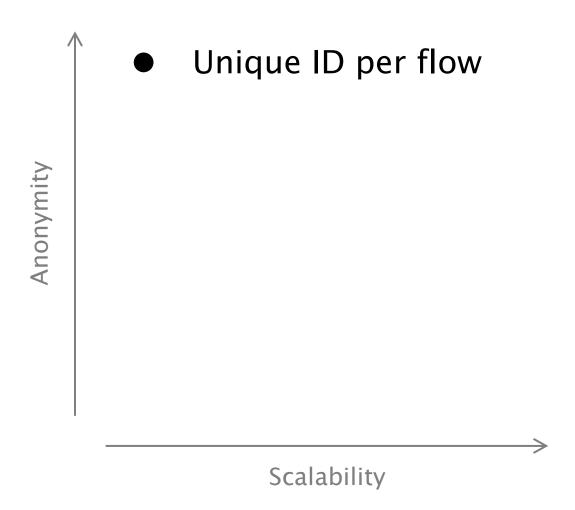


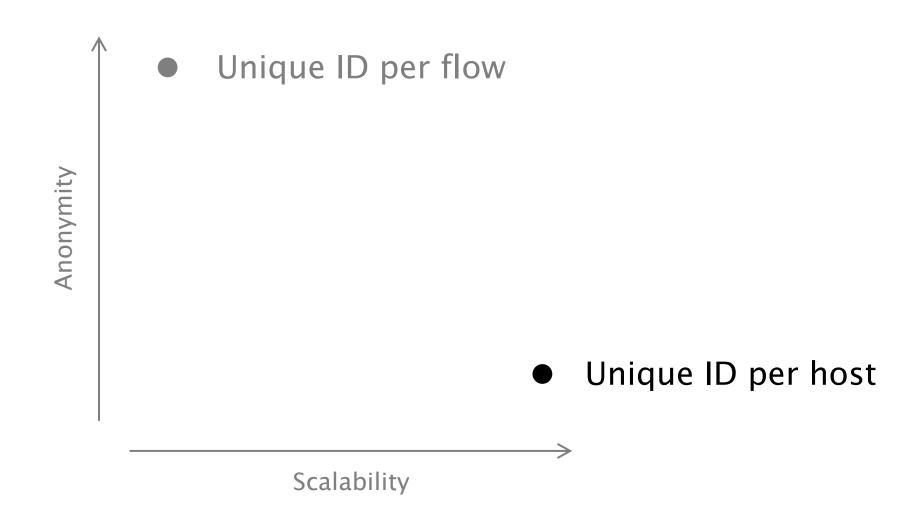
Measure information leakage & counteract attacker

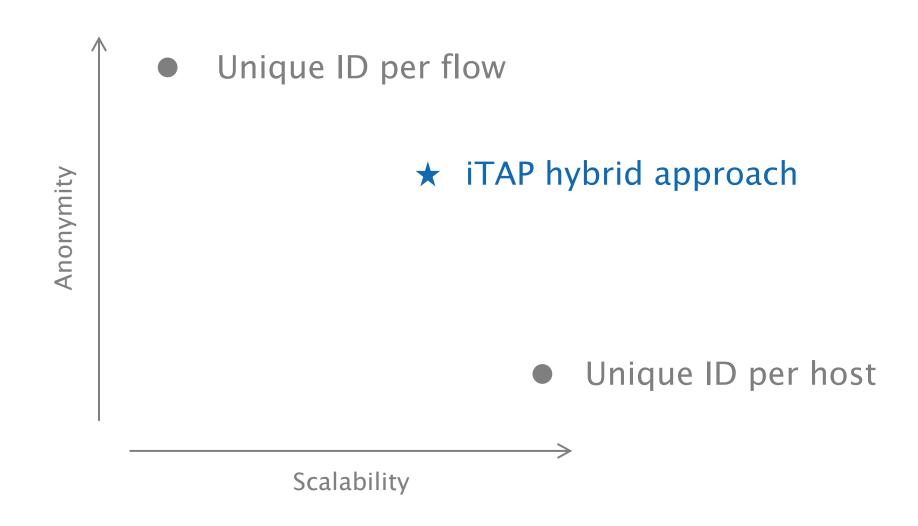


Solution for potential scalability-issues at Internet-facing edge



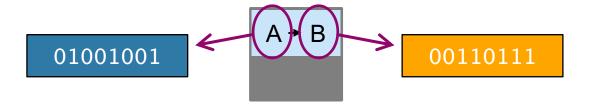




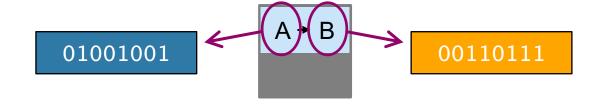




Map source and destination to IDs



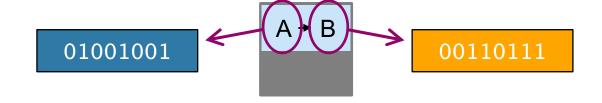
Map source and destination to IDs



Match-fields with arbitrary bitmasks

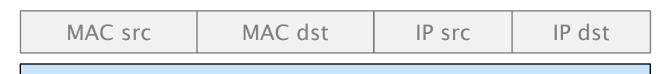
MAC src MAC dst	IP src	IP dst
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Map source and destination to IDs



Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits



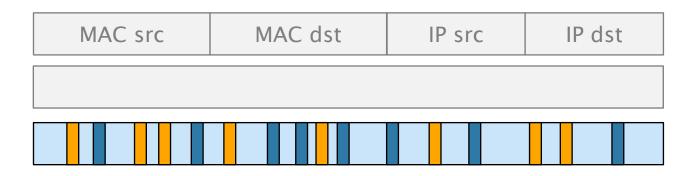
Map source and destination to IDs

01001001 (A) (B) 00110111

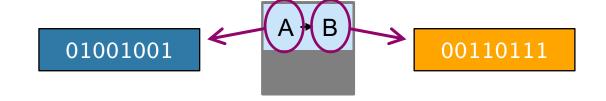
Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID



Map source and destination to IDs

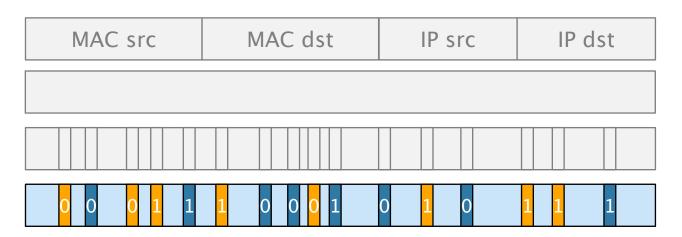


Match-fields with arbitrary bitmasks

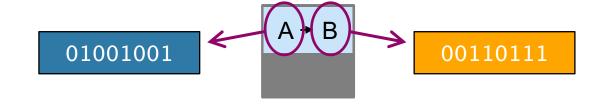
Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID

Add source and destination ID



Map source and destination to IDs



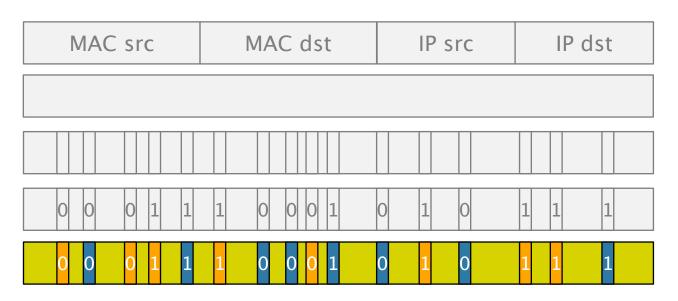
Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

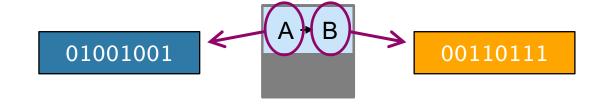
Randomly select bits that are used for source and destination ID

Add source and destination ID

Set other bits to random values



Map source and destination to IDs



Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID

Add source and destination ID

Set other bits to random values





Forwarding based on the destination ID

good scalability





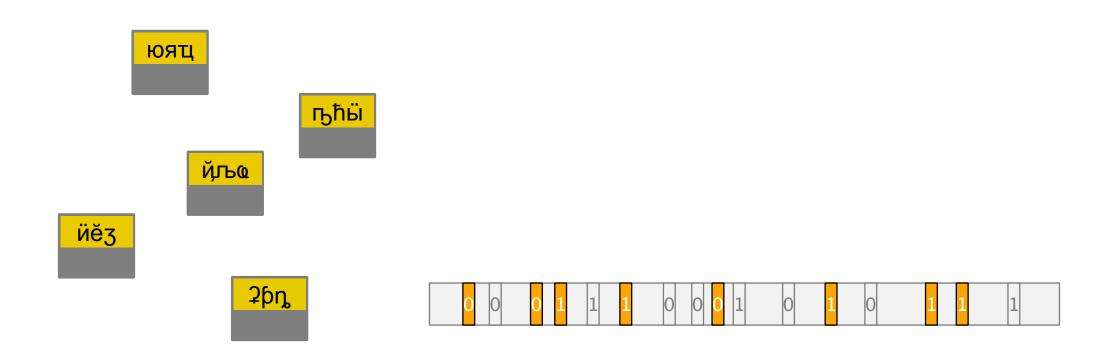
Eavesdropper cannot distinguish between random and non-random bits

→ good anonymity

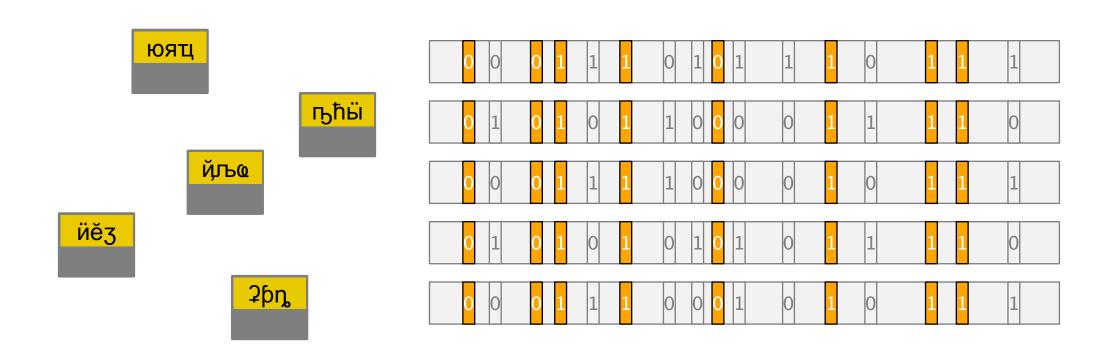




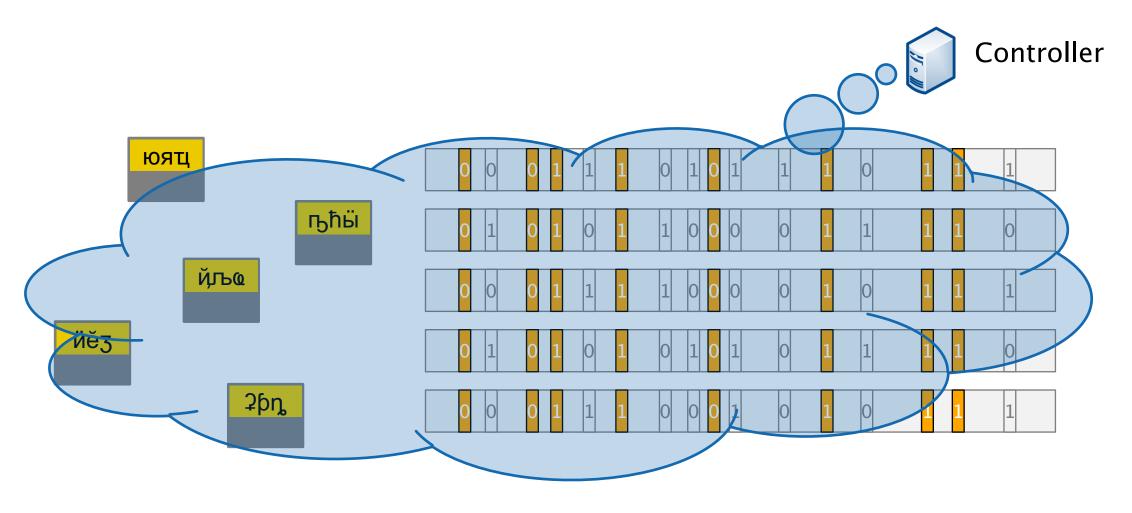
What if an attacker analyzes multiple flows?



What if an attacker analyzes multiple flows?

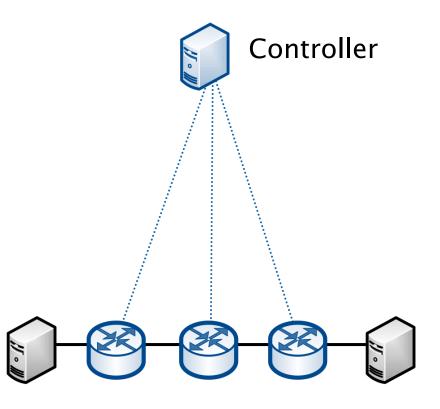


What if an attacker analyzes multiple flows?



The controller monitors the observed entropy for each link...

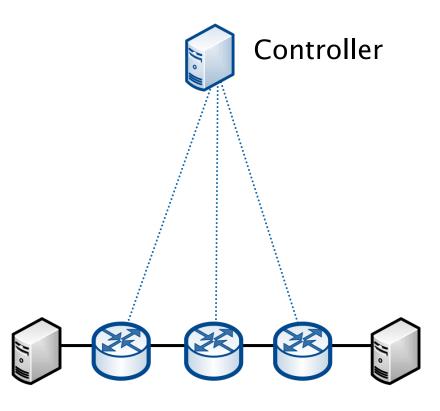
... and changes the encoding before an eavesdropper is able to break it.

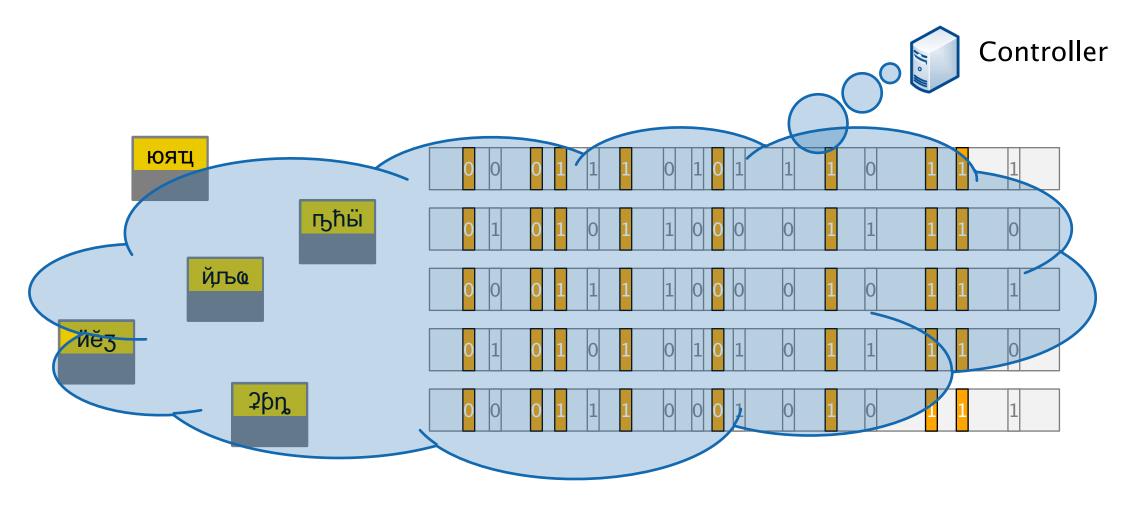


The controller monitors the observed entropy for each link...

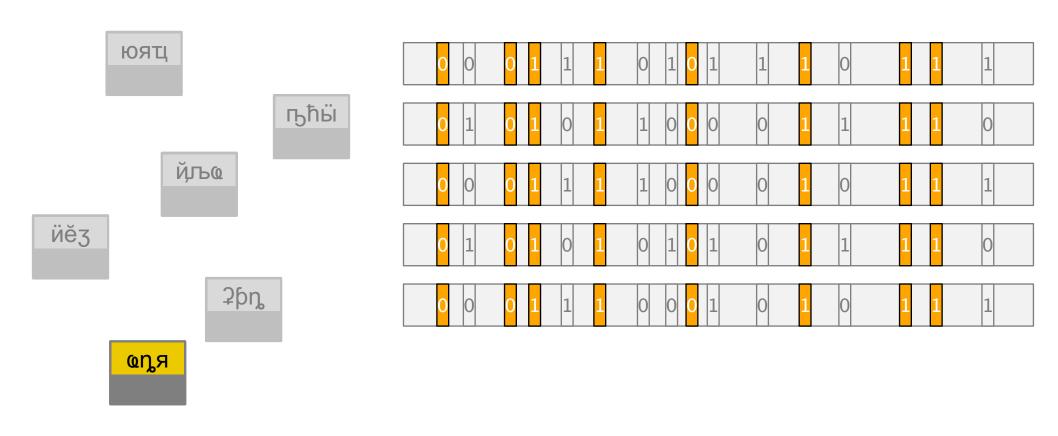
... and changes the encoding before an eavesdropper is able to break it.*

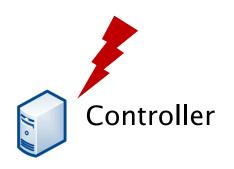
* According to the Unicity Distance

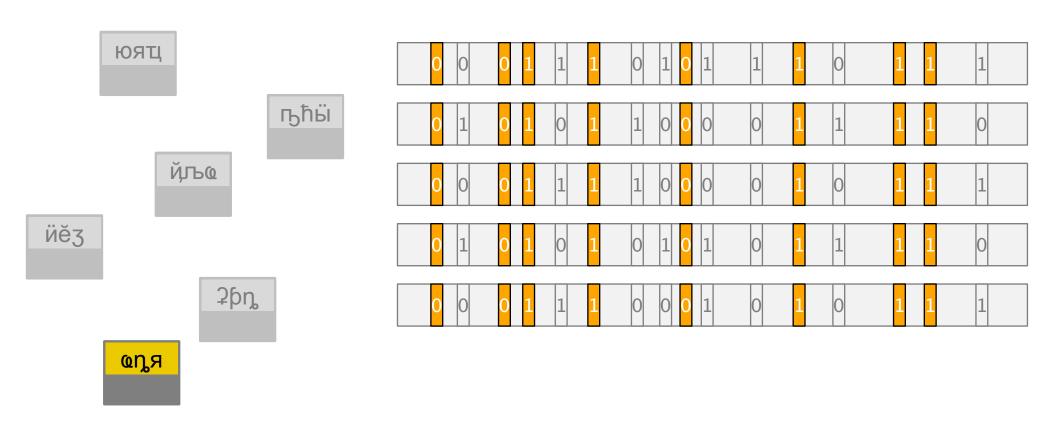














Forwarding based on the destination ID

good scalability





Forwarding based on the destination ID

- → good scalability
- → but what about the edge switches?





Distributed rewriting for better scalability at the network edge

E

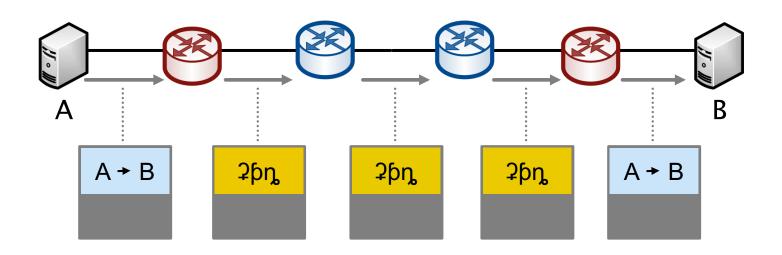
Forwarding

core switches

E

Rewriting

edge switches



Distributed rewriting for better scalability at the network edge



Forwarding

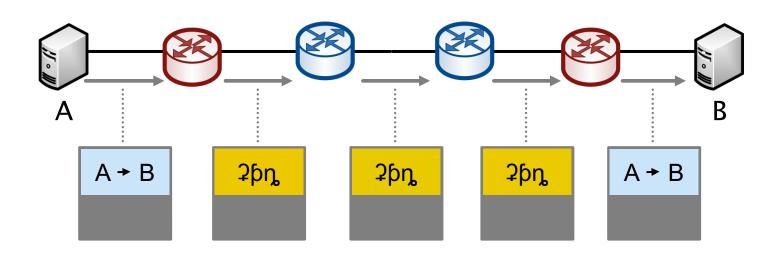
core switches

1 rule / destination (ID)



Rewriting

edge switches



Distributed rewriting for better scalability at the network edge



Forwarding

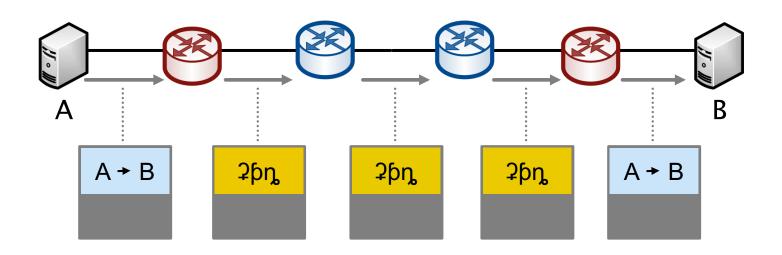
core switches

1 rule / destination (ID)



Rewriting

edge switches 1 rule / flow



Distributed rewriting for better scalability at the network edge



Forwarding

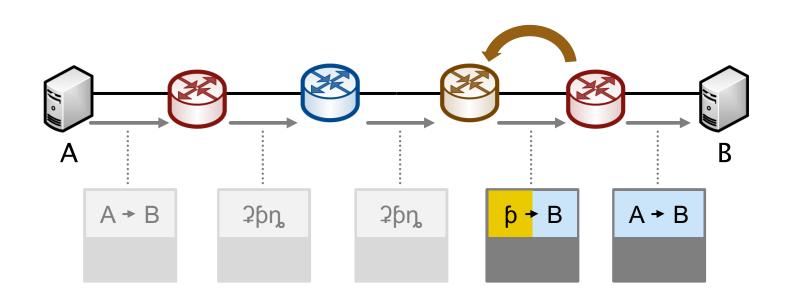
core switches

1 rule / destination (ID)



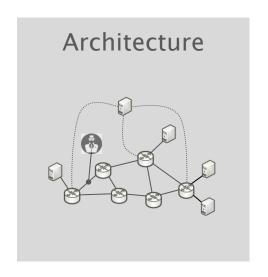
Rewriting

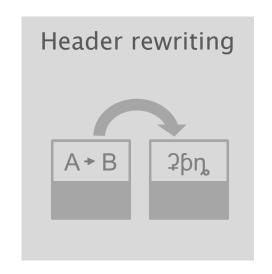
edge switches / 1 rule / flow core switches

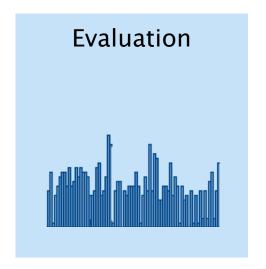


iTAP









iTAP evaluation based on real network traffic

iTAP evaluation based on

7 days of network traffic

400 hosts

128 M flows

iTAP evaluation based on real network traffic

7 days of network traffic 400 hosts 128 M flows

Indicators: controller actions / s

flow table updates / s

forwarding rules

iTAP works in practice

7 days of network traffic 400 hosts

128 M flows

avg max

200 700 controller actions / s

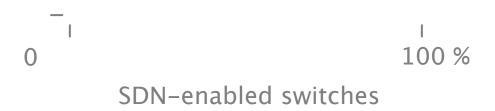
50 250 flow table updates / s

600 2.5 k forwarding rules

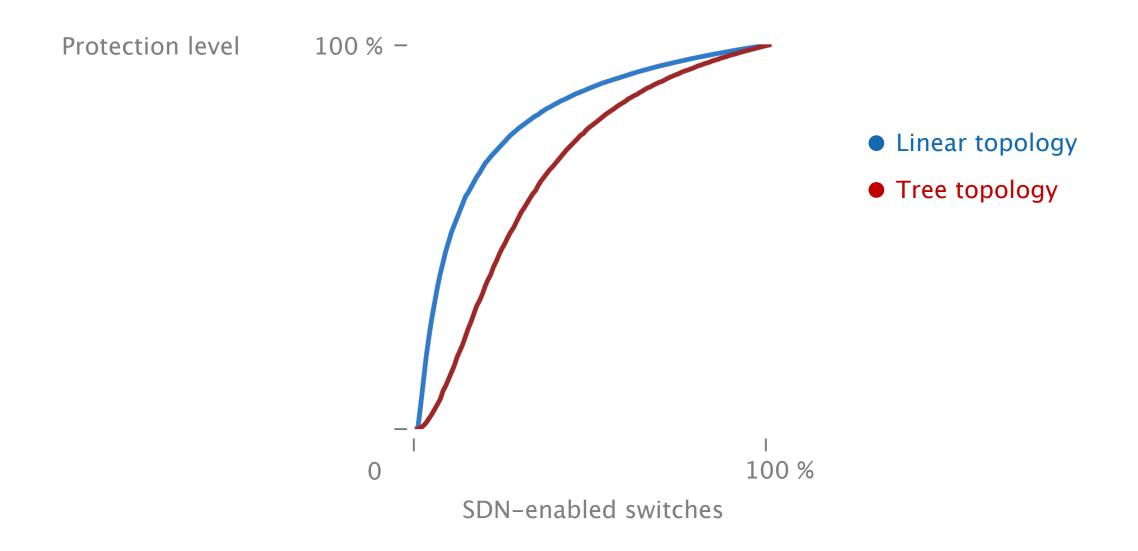
Only a small share of SDN switches is sufficient to protect a large share of the network traffic

Protection level 100 % -

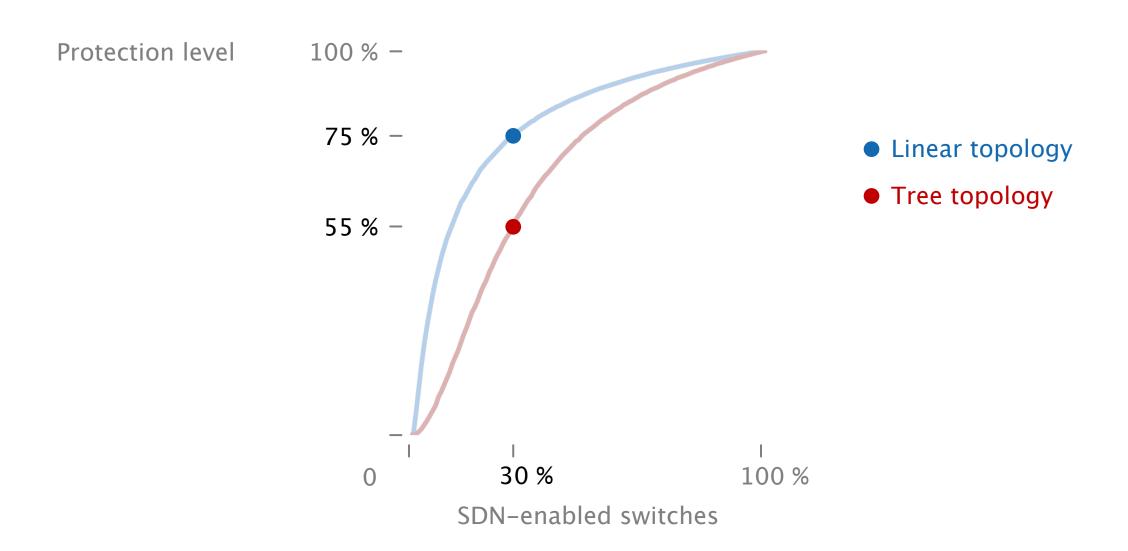
- Linear topology
- Tree topology



Only a small share of SDN switches is sufficient to protect a large share of the network traffic



Only a small share of SDN switches is sufficient to protect a large share of the network traffic



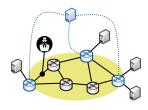
Outlook Improving network security through programmability



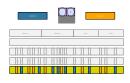
- Anonymity & privacy
- Detecting & locating attackers
- Deception techniques

Contributions

https://itap.ethz.ch



iTAP design



Scalable & anonymity-providing header rewriting scheme



iTAP prototype implementation



Evaluation based on real user traffic