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SDN-based Network Obfuscation

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N.S.A. May Have Hit Internet Companies at a Weak Spot The Internet companies' data centers are locked down with full-time security [...]. But between the data centers [...] information was unencrypted and an easier target for government intercept efforts, ... - The New York Times, Nov. 25, 2013

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GLOBAL

The Creepy, Long-Standing Practice of Undersea Cable Tapping

The newest NSA leaks reveal that governments are probing "the Internet's backbone." How does that work?



OLGA KHAZAN | JUL 16, 2013

The Washington Post Technology 4 6 Google encrypts data amid backlash against NSA spying By Craig Timberg September 6, 2013 💟 Google is racing to encrypt the torrents of information that flow among its data centers around the world in a bid to thwart snooping by the NSA and the intelligence agencies of foreign governments, company officials said Friday. The move by Google is among the most concrete signs yet that recent revelations about the National Security Agency's sweeping surveillance efforts have provoked significant backlash within an American technology industry that U.S. government officials long courted as a potential partner in Google's encryption initiative, initially approved last year, was accelerated in June as the tech giant struggled to guard its reputation as a reliable steward of user information amid controversy about the NSA's PRISM program, first reported in The Washington Post and the Guardian that month. PRISM obtains data from American technology companies, including Google, under various legal authorities.





[janitha.com]



This thesis vs. existing solutions

Alice

$ \frown $
Hi Bob,
$ \times $

source: Alice
destination: Bob
Hi Bob,

Payload encryption





This thesis vs. existing solutions





This thesis

SDN-based Network Obfuscation



This thesis

SDN-based Network Obfuscation

- Communication anonymity who is communicating with whom?
- Volume anonymity how much traffic flows between host X and Y?
- Topology anonymity how many hosts are in the network?





SDN-based Network Obfuscation

- Software-Defined Network New network architecture
- Network-based approach No modifications at end-hosts





Software-Defined Networking

Traditional network

closed software

closed hardware

[Cisco]





Software-Defined Networking

SDN

open

Traditional network



closed software

closed hardware





















With some SDN switches



Architecture





With some SDN switches

And a central controller



Architecture





With some SDN switches

And a central controller

Attacked by an eavesdropper



Architecture





With some SDN switches

And a central controller

Attacked by an eavesdropper

Protected by our system









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Packet from A to B enters the network





Ingress switch notifies controller





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Controller computes & installs flow rules





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Ingress switch obfuscates source and destination





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Core switch forwards obfuscated packet

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Egress switch de-obfuscates source and destination

How does the rewriting work?

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Rewriting as a trade-off between anonymity and scalability

Anonymity

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Anonymity

Unique ID per flow

Rewriting as a trade-off between

anonymity and scalability

Rewriting as a trade-off between anonymity and scalability

Unique ID per flow

Anonymity

Rewriting as a trade-off between anonymity and scalability

Anonymity

Unique ID per flow

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Rewriting scheme

Map source and destination to IDs

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Rewriting scheme

Map source and destination to IDs

Match-fields with arbitrary bitmasks

MAC src MAC dst IP src IP dst	MAC src	MAC dst	IP src	IP dst
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Map source and destination to IDs

Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

MAC src	MAC dst	IP src	IP dst

Map source and destination to IDs

Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID

MAC src	MAC dst	IP src	IP dst

Map source and destination to IDs

Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID

Add source and destination ID

MAC src	MAC dst	IP src	IP dst

Map source and destination to IDs

Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits

Randomly select bits that are used for source and destination ID

Add source and destination ID

Set other bits to random values

MAC src	MAC dst	IP src	IP dst
0 0 0 1 1	1 0 0 0 1	0 1 0	1 1 1
0 0 0 1 1	1 0 0 0 1	0 1 0	1 1 1

Map source and destination to IDs

Match-fields with arbitrary bitmasks

Interpret as bit-string of 160 bits Randomly select bits that are used for

source and destination ID

Add source and destination ID

Set other bits to random values

MAC src	MAC dst	IP src	IP dst
0 0 0 1 1	1 0 0 0 1	0 1 0	1 1 1
0 0 0 1 1	1 0 0 0 1	0 1 0	1 1 1

Obfuscation controller compared with Floodlight in default configuration

Resource usage in switches # flow table entries

Switch load # flow table updates / s

Controller load # flows / s

Network performance RTT and bandwidth

Evaluation

Obfuscation controller compared with Floodlight in default configuration

Resource usage in switches # flow table entries

Switch load # flow table updates / s

Controller load # flows / s

Network performance RTT and bandwidth

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Follow-up work

iTAP: In-network Traffic Analysis Prevention using Software-Defined Networks

https://itap.ethz.ch

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ABSTRACT

Advances in layer 2 networking technologies have fostered the deployment of large, geographically distributed LANs, Due to their large diameter, such LANs provide many vantage points for wiretapping. As an example, Google's internal network was reportedly tapped by governmental agencies, forcing the Web giant to encrypt its internal traffic. While using encryption certainly helps, eavesdroppers can still access traffic metadata which often reveals sensitive information, such as who communicates with whom and which are the critical hubs in the infrastructure.

This paper presents iTAP, a system for providing strong anonymity guarantees within a network. iTAP is networkbased and can be partially deployed. Akin to onion routing, iTAP rewrites packet headers at the network edges by leveraging SDN devices. As large LANs can see millions of flows, the key challenge is to rewrite headers in a way that guarantees strong anonymity while, at the same time, scaling the control-plane (number of events) and the dataplane (number of flow rules), iTAP addresses these challenges by adopting a hybrid rewriting scheme. Specifically, iTAP scales by reusing rewriting rules across distinct flows and by distributing them on multiple switches. As reusing headers leaks information, iTAP monitors this leakage and adapts the rewriting rules before any eavesdropper could provably de-anonymize any host.

We implemented iTAP and evaluated it using real network traffic traces. We show that iTAP works in practice, on existing hardware, and that deploying few SDN switches is enough to protect a large share of the network traffic.

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CCS Concepts

 Security and privacy → Pseudonymity, anonymity and untraceability; Network security:
●Networks → Network privacy and anonymity; Programmable networks;

Keywords

Anonymous communication; wiretapping; SDN

1. INTRODUCTION

Since the Snowden revelations, it is well-known that network eavesdropping was (and probably still is) performed in the Internet core, particularly on undersea cables [22]. While worrying, these threats can be mitigated to a large degree by hiding connection metadata (e.g., using Tor [14]) and by relying on pervasive encryption (e.g., using VPNs).

However, network eavesdropping is not limited to the Internet backbone. As enterprise networks become bigger in terms of users and physical reach, they, too, become susceptible to wiretapping. Indeed the advent of new layer 2 technologies such as TRILL [7], Shortest Path Bridging [30], or SDN-based solutions [34] enables network administrators to build large LAN zones that can easily span several thousand devices and users. Due to their large physical diameter, such networks inevitably exhibit many vantage points for wiretapping. Actually, the majority of the attacks is now performed by insiders, i.e. malicious insiders or inadvertent actors [4] that act from within the network, rather than by remote attackers. As an example, Google's Wide-Area Network (WAN) - the private network connecting Google's data centers - was reportedly tapped by governmental agencies, forcing the Web giant to encrypt its internal traffic [36].

While encrypting internal traffic protects the payload of connections, an in-network attacker can still monitor and analyze the unencrypted packet headers, i.e., the metadata, In particular, the MAC addresses and, in case of SSL/TLS application laver encryption, also the IP addresses of the source and destination hosts along with the source and destination ports used for the communication.

By analyzing these unencrypted header fields, an attacker can gain useful information about: i) which hosts communicate; ii) the topology of the network; iii) the addresses of

This thesis

+ Partial deployment

- + Improved scalability at network edge
- + Evaluation based on real user traffic

https://itap.ethz.ch

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Contributions

https://itap.ethz.ch

Network-based design

Scalable & anonymity-providing header rewriting scheme

Prototype implementation (open source)

Evaluation

Contributions

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Network-based design

Scalable & anonymity-providing header rewriting scheme

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Prototype implementation (open source)

and added

Evaluation

Thank you! Questions?